

RUTGERS Optimizing LT Codes for Multicast in Heterogeneous Networks†

WINLAB | Wireless Information Network Laboratory

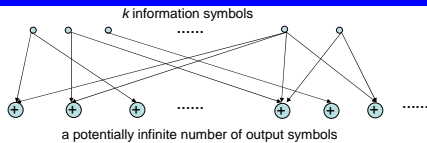
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Motivation

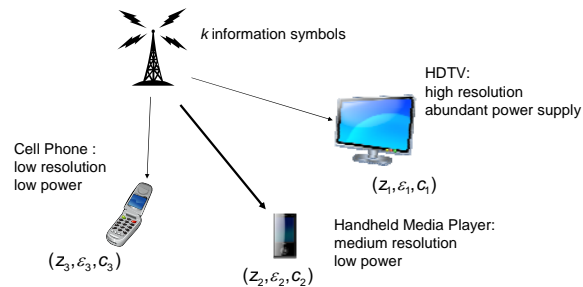
- Optimize LT rateless codes for multicast in a heterogeneous network of
 - Diverse Source-Sink Channel Conditions**
 - Non-Uniform Demand**
 - Both Decoding and Non-Decoding Sinks**

Degree Distribution of LT Codes



- Degree of an output symbol = the number of information symbols selected to be linearly combined.
- $p_i^{(k)} = \text{Prob}[\text{output symbol is of degree } i], i=1,2,\dots,k$
- Conventionally the “ideal soliton distribution” or its variations are used.
- Suboptimal iterative “belief propagation” (BP) decoding

System Model: A Heterogeneous Network



- k : the total number of information symbols
- Degree distribution of LT symbols encoded by the source $P^{(k)}(x) = p_1^{(k)}x + p_2^{(k)}x^2 + \dots + p_k^{(k)}x^k$
- Sink clusters characterized by (z_i, c_i, ϵ_i) , size $n_i, i=1,2,\dots,l$
 - z_i : the fraction of information symbols the sink intends to recover
 - $c_i = 1$ for a BP decoding and $= 0$ for a non-decoding sink
 - ϵ_i : erasure rate of the BEC link connecting the source to the sink
 - t_i : latency, the number of output symbols transmitted over the link until the demand of the sink is satisfied, normalized by k

Optimization in the Asymptotic Regime

Performance Measures

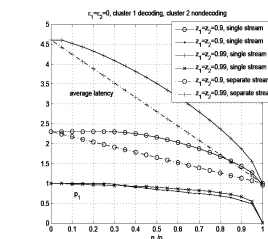
- Min-Max Latency:** minimize $\rho, \max_i t_i$
 - minimize the total latency at the source side
- Max-Min Throughput:** maximize $\rho, \min_i z_i/t_i$
 - Throughput measures the actual channel degradation and provides reference for the pricing of services.
- Max-Min Channel Utilization:** maximize $\rho, \min_i z_i/(t_i(1-\epsilon_i))$
 - Expect worse service on worse channels
- Min Average Latency** minimize $\rho, t_i, \frac{\sum n_i t_i}{\sum n_i}$
 - Optimize the average performance of all sink nodes

Optimization Constraints

- Ripple**
 - the set of output symbols of (induced) degree-1 during BP decoding of LT codes
 - The decoding process halts when the ripple becomes empty
- Optimization constraints:**
 - keep ripple size > 0 until the unrecovered fraction goes down to $1-z$
 - For a decoding sink cluster $i: t_i(1-\epsilon_i)P'(1-u) + \ln u > 0, 1-z_i < u \leq 1$
 - For a non-decoding sink cluster $i: t_i(1-\epsilon_i)p_i + \ln u > 0, 1-z_i < u \leq 1$
- Derivation of the Optimization Constraints
 - Theorem**[Maatouk et al.'09] (the expected ripple size (rk) as a function of the number of unrecovered information symbols (uk) and the number of output symbols collected (wk))

$$r(u) = wu(P^{(k)}(1-u) + \frac{1}{w} \ln u) + O\left(\frac{1}{k}\right)$$
 - As $k \rightarrow \infty, P^{(k)}(u) \rightarrow P(u), r(u) = wu(P(1-u) + \frac{1}{w} \ln u)$

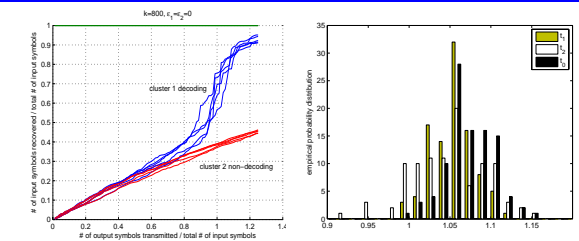
Minimizing Average Latency



Minimum average latency and achieving p_i versus the fraction of decoding sinks, $z_i = Z_2$.

- Perfect channels
- 2 clusters, one decoding, one non-decoding
- Uniform demand
 - Minimum average latency drops as the fraction of decoding sinks grows larger
 - p_i of the optimized degree distribution falls slowly as the fraction of decoding sinks grows larger

Finite-Length Simulations



Decodable fraction v.s. # of output symbols collected; $z_1=0.8, z_2=0.4$, perfect channel, $k=800$; 5 sample runs

Histograms for t_1, t_2 and $t_0 = \max\{t_1, t_2\}$; 100 sample runs, mean of t_0 is 1.0718, standard deviation 0.0300; optimization result $t_0 = 1.0473$

Source Latency Comparison

- Total number of output symbols transmitted by the source

	(z_1, c_1, ϵ_1) (z_2, c_2, ϵ_2)	Scheme A0	Scheme A1	Scheme A2	Scheme A12
cluster 1	(0.98, 1.0)	1.5636	0.9914	3.9120	0.9914
cluster 2	(0.72, 0.0)	1.4106	∞	1.2730	1.2730
cluster 1	(0.98, 1.0)	1.6220	0.9914	1.9959	0.9914
cluster 2	(0.63, 1.0, 0.5)	1.6220	1.9828	1.5782	1.5782

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Schemes

- A0-A2: the source sends a **single stream** to all sinks
- A0: minimize the **max latency**.
- A1: minimize the **latency of cluster 1**.
- A2: minimize the **latency of cluster 2**.
- A12: The source sends **two independent streams** to the clusters, each minimizing the **latency of the targeted cluster**.

Conclusions and Future Work

- We have optimized LT codes for multicast in a highly heterogeneous network
- Optimization is through the design of the code degree distribution
- Optimized performance shows **significant savings in the total bandwidth consumption** compared to sending separately encoded streams to different types of sink nodes
- Finite-length simulations **confirm the usability of the asymptotic results to guide the finite-length code design**
- Future Work: More complex network topologies
 - Cooperation between sinks?
 - Multi-hop relay?

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