

On the Capacity of Deterministic Interference Channels Beyond Two Users

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Interference Channels and Deterministic Models

Capacity of a Gaussian Interference Channel

- A long standing open problem
- Recent Advancements made by using the 'deterministic model'
- Capacity of the canonical 2 user IC has been determined upto 1 bit [Etkin, Tse, Wang '07]

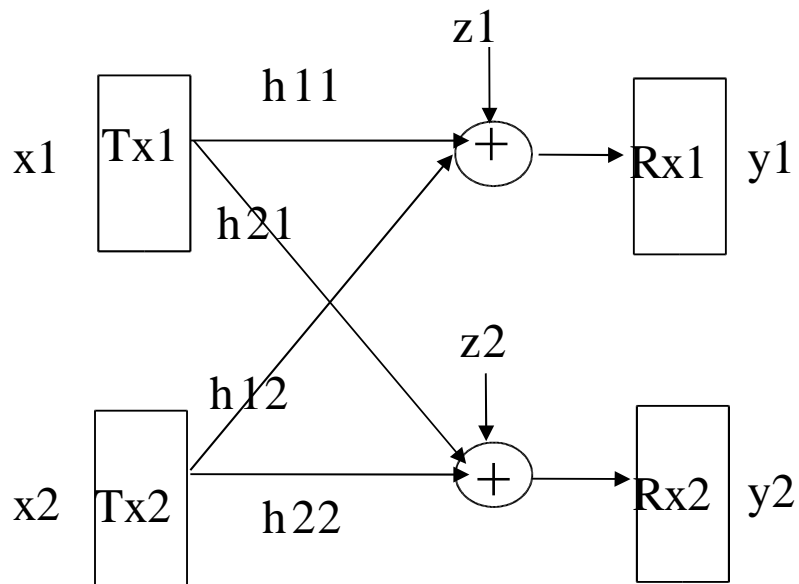


Fig 1. Two User Gaussian Interference Channel

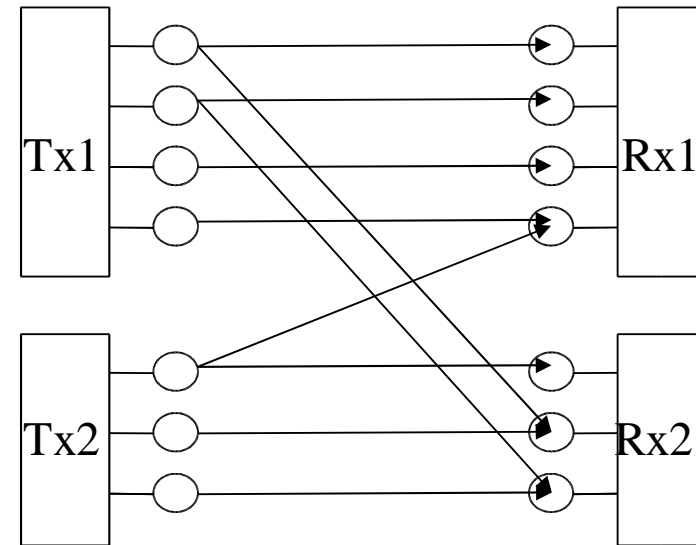


Fig 2. Two User Deterministic Interference Channel

Deterministic Model

- First proposed in the context of Gaussian Relays [Avestimehr, Diggavi, Tse '07]
- Naturally models the Gaussian Channel
- Simpler to analyze
- Provides intuition towards 'capacity achieving' schemes in the Gaussian Case

Known Capacity Results for Deterministic IC

Capacity In 2 User Case

- Known for a General Class of Deterministic IC due to [*El-Gamal, Costa, '82*]
- Alternatively derived for a more special case by [*Bresler, Tse '08*]
- Can be achieved by the Han-Kobayashi schemes [*Han, Kobayashi '81*],

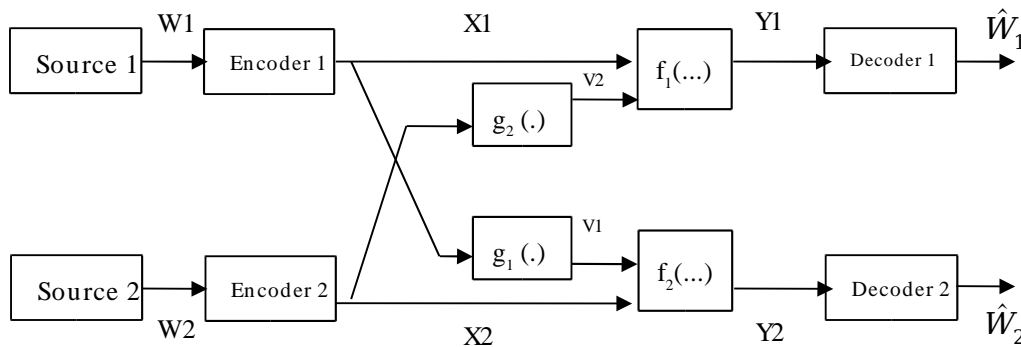


Fig 3. A General Deterministic IC

$$\begin{aligned}
 R_1 &\leq H(Y_1 | V_2) \\
 R_2 &\leq H(Y_2 | V_1) \\
 R_1 + R_2 &\leq H(Y_1 | V_1 V_2) + H(Y_2) \\
 R_1 + R_2 &\leq H(Y_2) + H(Y_2 | V_1 V_2) \\
 R_1 + R_2 &\leq H(Y_1 | V_1) + H(Y_2 | V_2) \\
 2R_1 + R_2 &\leq H(Y_1) + H(Y_1 | V_1 V_2) + H(Y_2 | V_2) \\
 R_1 + 2R_2 &\leq H(Y_1 | V_1) + H(Y_2) + H(Y_2 | V_1 V_2)
 \end{aligned}$$

Capacity Region for the Model in Figure 3.

Beyond 2 Users

- Capacity is known for Many-to-One and One-to-Many Channels [*Bresler, Parekh, Tse, 08*]
- Three User Lattice Coding strategies for Gaussian Case [*Sridharan, Jafri, Viswanath, Jafar, Shamai '08*]

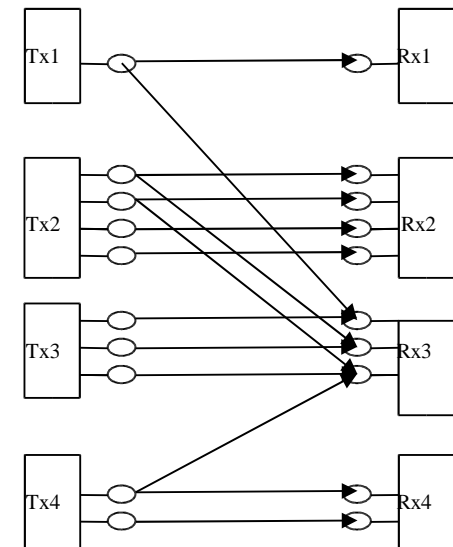
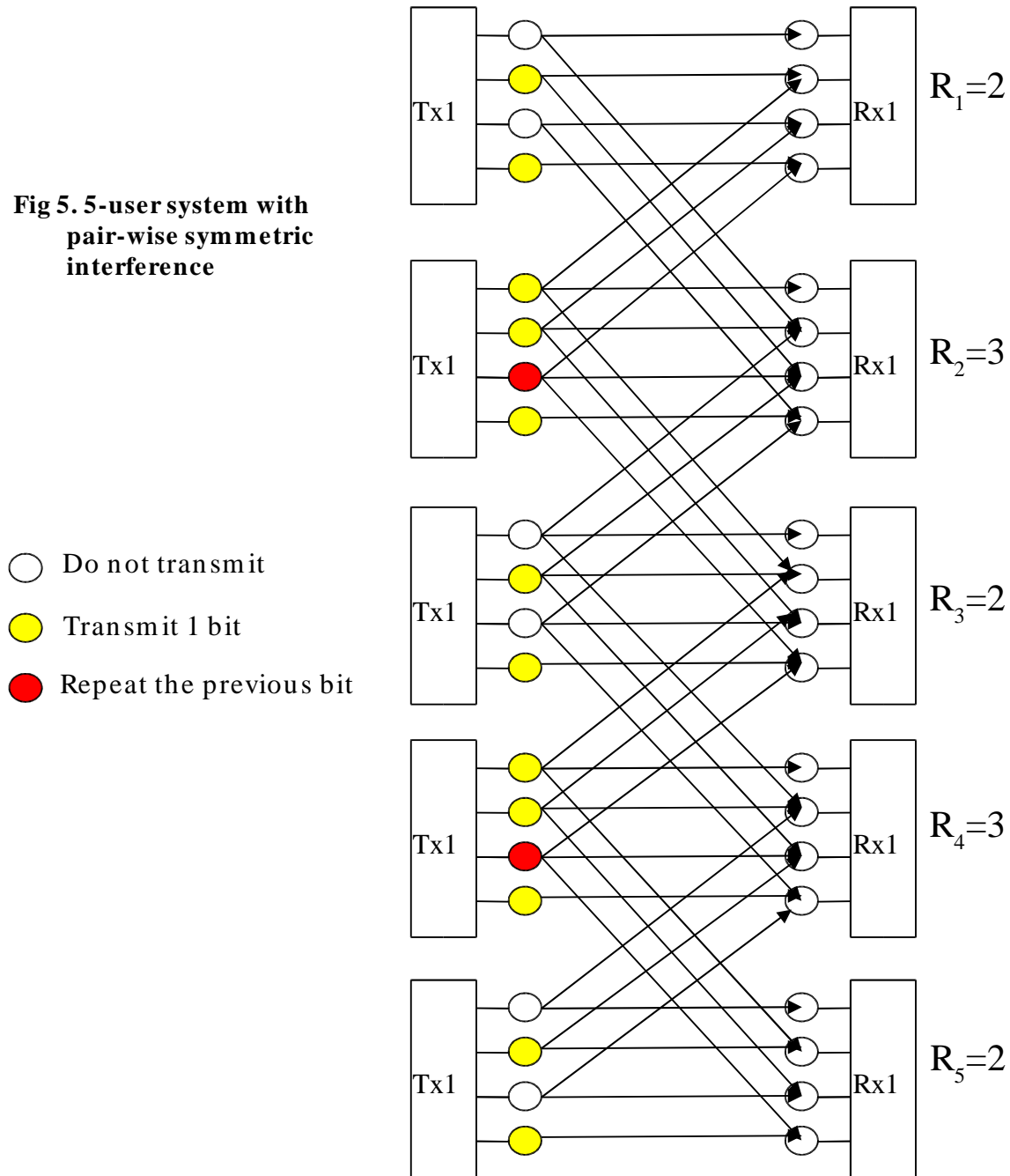


Fig 4. Many -to - One Deterministic IC

New Capacity Results for 'n- user' Pairwise Symmetric Interference

Fig 5. 5-user system with pair-wise symmetric interference



Lemma: For an 'n' user symmetric pairwise interference channel, the sum capacity can be achieved by directly extending the two-user symmetric interference case.

Key Observation : Considering three adjacent users at a time, due to the symmetry in the interference pattern, the achievable two-user rates are also achievable.

In this example, we have $n_{11} = n_{22} = n_{33} = n_{44} = n_{55} = 4$ while the symmetry conditions are given by $n_{12} = n_{21} = n_{23} = n_{32} = n_{34} = n_{43} = n_{45} = n_{54} = 3$

$R_1 = 2, R_2 = 3$ achieves capacity for a two user system with this interference pattern [Berry, Tse '08]

Three User Pairwise Interference System

Motivation

- simplest extension beyond 2 user system motivated by the previous example.
- one user only gets interfered by its two nearest users and thus is just one step beyond a two user system that captures the peculiarity of an 'n' user system
- models the interference pattern of cellular base stations in a linear, say highway, network.

A General 3 User Deterministic Interference Model

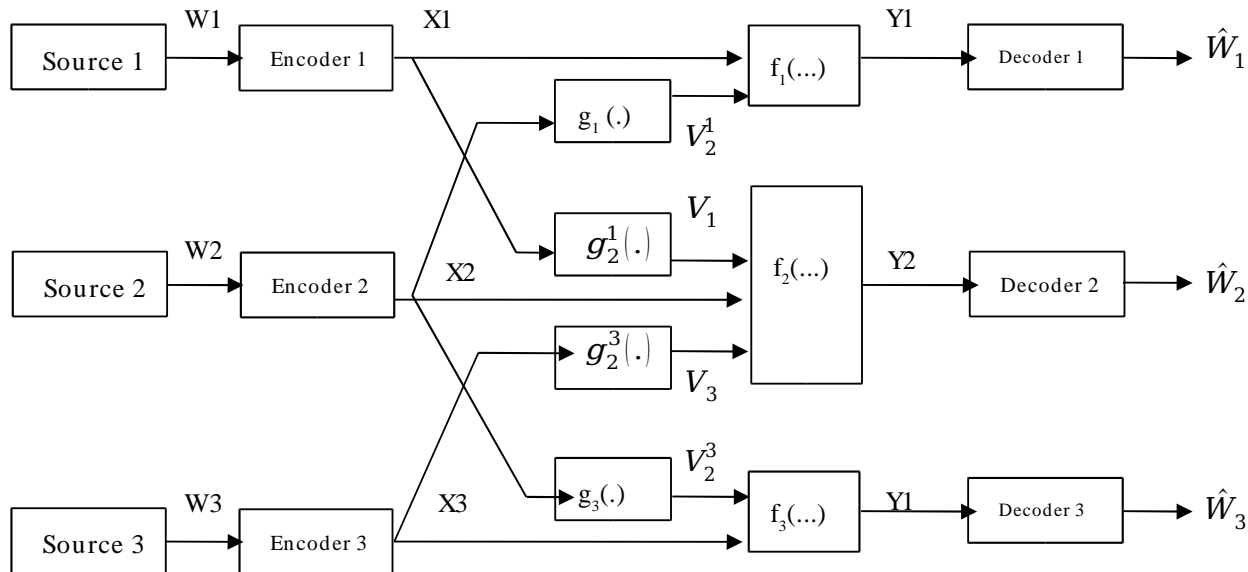


Fig 6. A General 3 user pairwise interfering Deterministic IC

Sum Rate Upper Bounds

Bounds

$$\left. \begin{aligned} (1) \quad R_1 &\leq H(Y_1 | V_2) \\ (2) \quad R_2 &\leq H(Y_2 | V_1) \\ (3) \quad R_3 &\leq H(Y_3 | V_1) \end{aligned} \right\} \text{Individual rate bounds}$$

$$(4) \quad R_1 + R_2 + R_3 \leq .5 * (H(Y_1) + H(Y_1 | V_1 V_2^1) + H(Y_2 | V_2^1) + H(Y_3) + H(Y_3 | V_3 V_2^3) + H(Y_3 | V_2^3)) \quad \left. \vphantom{R_1 + R_2 + R_3} \right\} \text{Comes from Two-user Sum Rate Bounds}$$

$$(5) \quad R_1 + R_2 + R_3 \leq H(Y_1 | V_1 V_2^1) + H(Y_2) + H(Y_3 | V_3 V_2^3) + H(V_1 V_3) - H(g(V_1 V_3))$$

$$(6) \quad R_1 + R_2 + R_3 \leq H(Y_1) + H(Y_2 | V_1 V_3 V_2^1 V_2^3) + H(Y_3) - I(V_2^1; V_2^3)$$

$$(7) \quad R_1 + R_2 + R_3 \leq H(Y_1 | V_1) + H(Y_2 | V_2^1 V_2^3) + H(Y_3 | V_3) + H(V_1 V_3) - H(g(V_1 V_3))$$

$$(8) \quad R_1 + R_2 + R_3 \leq H(Y_1 | V_1 V_2^1) + H(Y_2 | V_1 V_3 V_2^1 V_2^3) + H(Y_3) + H(V_1) + H(V_2^1) + H(V_2^1 | V_2^3) - H(V_2^3)$$

$$(9) \quad R_1 + R_2 + R_3 \leq H(Y_1) + H(Y_2 | V_1 V_3 V_2^1 V_2^3) + H(Y_3 | V_3 V_2^3) + H(V_3) + H(V_2^3) + H(V_2^3 | V_2^1) - H(V_2^1)$$

$$(10) \quad R_1 + R_2 + R_3 \leq H(Y_1 | V_1 V_2^1) + H(Y_2) + H(Y_3) + H(V_1) - H(V_2^3) - H(g(V_1 V_3))$$

$$(11) \quad R_1 + R_2 + R_3 \leq H(Y_1) + H(Y_2) + H(Y_3 | V_3 V_2^3) + H(V_3) - H(V_2^1) - H(g(V_1 V_3))$$

} Three-user Sum Rate Bounds

Observations

1. Sum rate bounds for this model differs from the two user model in terms of the presence of the function $g(\cdot)$ that models the interaction between the interferences from users 1 and 3 at the receiver of user 2.

2. Unlike the 2 user case, uniform input distributions may not maximize the upper bounds in the RHS of all the above inequalities

Scenarios where bounds can be tight - I

Low Interference Scenario

$$n_{11} > n_{12} + n_{21}, \quad n_{22} > \max(n_{21}, n_{23}) + \max(n_{12}, n_{32}),$$

$$n_{33} > n_{23} + n_{32},$$

Assuming $n_{21} > n_{23}, n_{12} > n_{32}$
the bound (7) can be reformulated as

$$R_1 + R_2 + R_3 \leq \max(n_{12}, (n_{11} - n_{21})^+) + \max(n_{32}, (n_{33} - n_{23})^+) + \max(\max(n_{21}, n_{23}), (n_{22} - \max(n_{12}, n_{32}))^+) + \min(n_{21}, n_{23})$$

$$= n_{11} - n_{21} + n_{33} - n_{23} + n_{22} - n_{12} + n_{23}$$

$$= n_{11} - n_{21} + n_{33} + n_{22} - n_{12}$$

An achievable scheme: Transmit from all the levels at Tx 1 and Tx3. Transmit from the free levels at Tx 2 which are neither interfered with nor do interfere. This gives a sum rate of

$$n_{11} + n_{33} + (n_{22} - \max(n_{12}, n_{32}) - \max(n_{21}, n_{23}))^+$$

$$= n_{11} + n_{33} + n_{22} - n_{12} - n_{21}$$

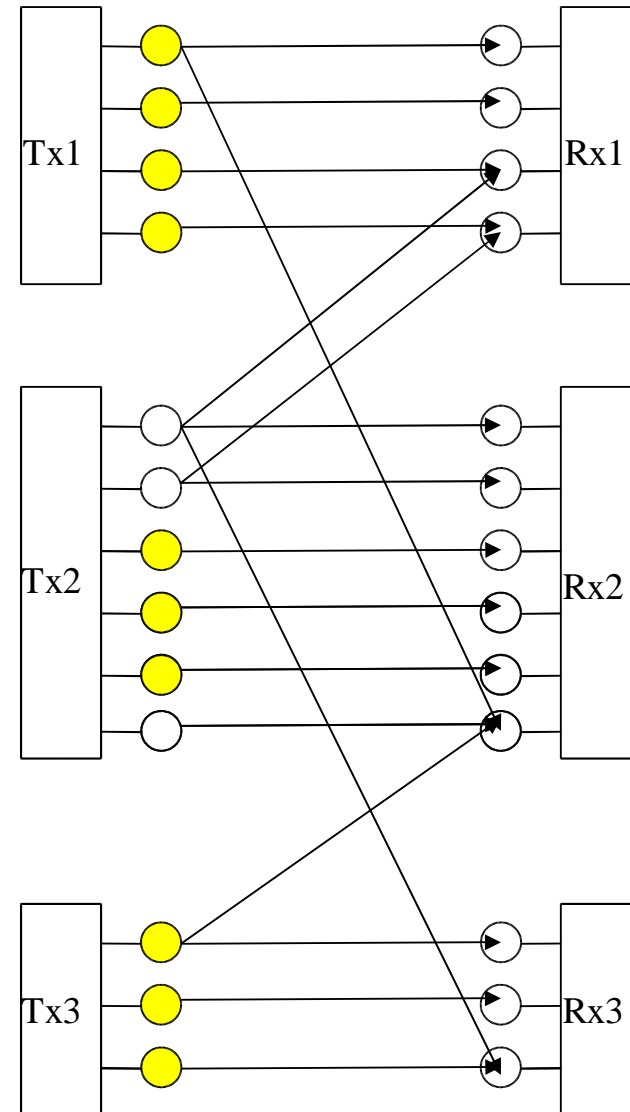
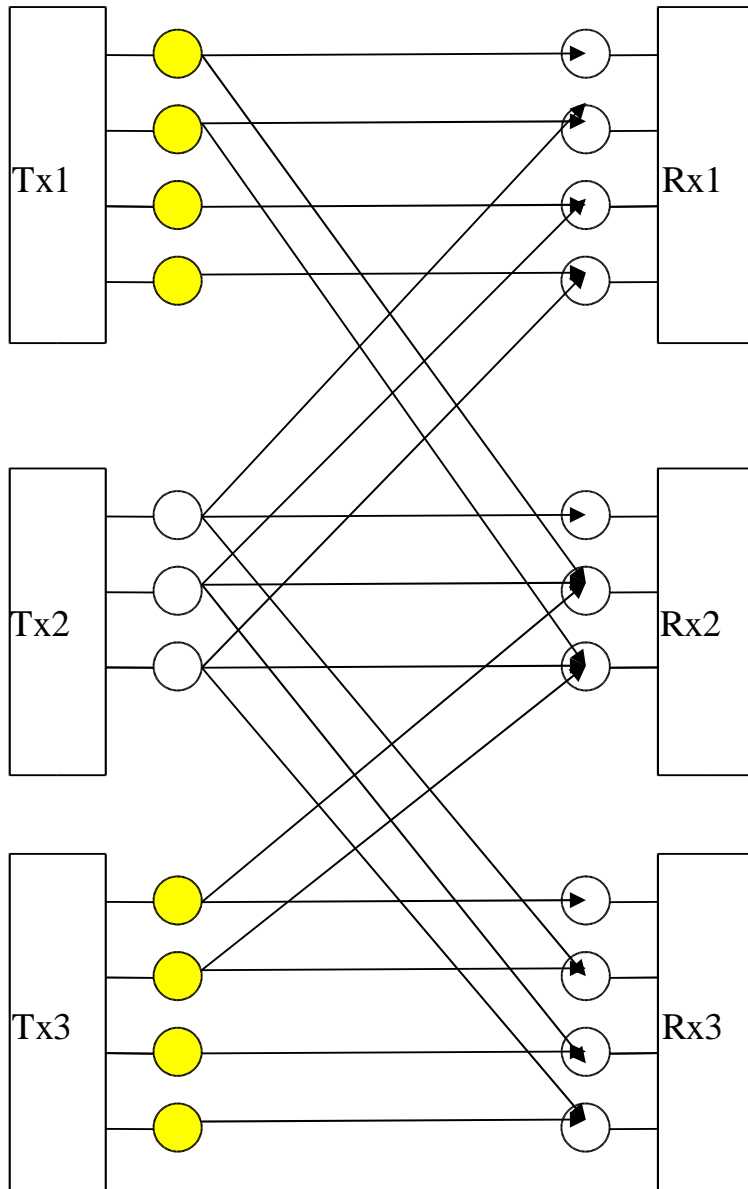


Fig 7. Low Interference Case where we can determine the sum capacity

Scenarios where bounds can be tight - II



A some-what Symmetric Scenario

The Interference relation between user 1 and user 2 is the same as that between user 3 and user 2, i.e.,

$$n_{12} = n_{32}, \quad n_{21} = n_{23}, \quad n_{11} = n_{33}$$

If such a symmetry exists, either the upper bound (4) or the upper bound derived from 2 user sum-rates is always tight.

In this example, upper bound (4) is tight as

$$\begin{aligned} 2*(R_1 + R_2 + R_3) &\leq \max(n_{11}, n_{12}) + (n_{11} - n_{21})^+ + \max(n_{21}, (n_{22} - n_{12})^+) \\ &\quad + \max(n_{33}, n_{32}) + (n_{33} - n_{23})^+ + \max(n_{23}, (n_{22} - n_{32})^+) \\ \Rightarrow R_1 + R_2 + R_3 &\leq .5*(n_{11} + n_{11} - n_{21} + n_{21} + n_{33} + n_{33} - n_{23} + n_{23}) \\ &= n_{11} + n_{33} \end{aligned}$$

An achievable strategy is to transmit from all levels of the users 1 and 3 while user 2 shuts off

Fig 8. A some-what Symmetric Interference Case where we can determine the sum capacity

What's Next and References

We plan to

- Characterize all the cases where the sum rate upperbounds are tight
- Extend our results to n-users
- Look at the Game Theoretic aspects of the 3 and then n user Interference games on the deterministic channel
- Use all the above results for extending to the Gaussian counterparts

References

1. Gaussian Interference Channel Capacity to Within One Bit [*Etkin, Tse, Wang '07*]
2. The Capacity Region of A Class of Deterministic Interference Channel [*El-Gamal, Costa, '82*]
3. The Approximate Capacity of the Many-to-One and One-to-Many Gaussian Interference Channels [*Bresler, Parekh, Tse, 08*]
4. Information Theoretic Games on Interference Channels [*Berry, Tse '08*]
5. A Deterministic Approach to Wireless Relay Networks [*Avestimehr, Diggavi, Tse '07*]
6. A New Achievable Rate Region for the Interference Channel [*Han, Kobayashi '81*]
7. A Layered Lattice Coding Scheme for a Class of Three User Gaussian Interference Channels [*Sridharan, Jafriani, Viswanath, Jafar, Shamai '08*]
8. The Two-user Interference Channel: A Deterministic View [*Bresler, Tse '08*]