

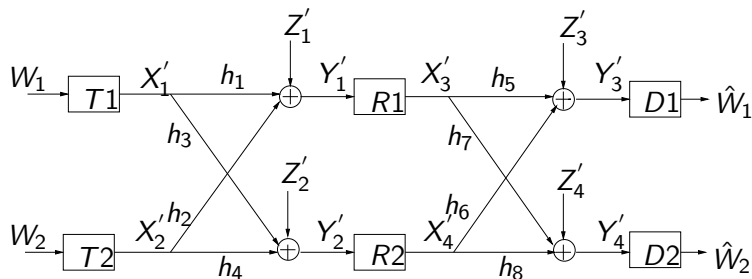
Two Hop Interference Networks

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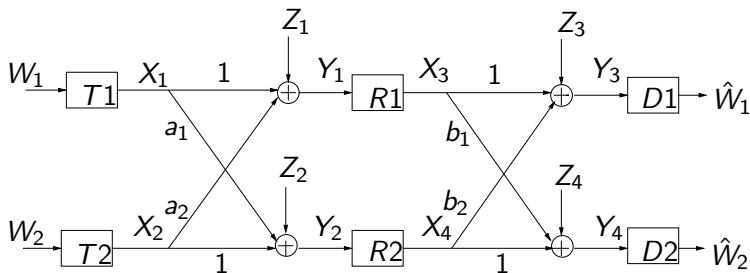
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General Channel Model and Related Work



- O. Simeone, O. Somekh, Y. Bar-Ness, H. V. Poor, and S. Shamai, 2007 (explored the coherent combining gain of decode-and-forward scheme)
- C. Thejaswi, A. Bennatan, J. Zhang, R. Calderbank, and D. Cochran, 2008 (focused on the second hop and explored the possibilities of utilizing the common messages from the first hop)

Standard Form



Consider Symmetric Channel Parameters

$$a_1 = a_2 \triangleq a$$

$$b_1 = b_2 \triangleq b$$

$$P_{11} = P_{12} \triangleq P_1$$

$$P_{21} = P_{22} \triangleq P_2$$

Case I: $0 < a < 1, 0 < b < 1$

- The first hop: Han and Kobayashi's Scheme

$$R_p^{(1)} = \gamma \left(\frac{\alpha P_1}{1 + a^2 \alpha P_1} \right)$$

$$R_c^{(1)} = \min \left\{ \gamma \left(\frac{a^2 \bar{\alpha} P_1}{\sigma_1^2} \right), \frac{1}{2} \gamma \left(\frac{(1 + a^2) \bar{\alpha} P_1}{\sigma_1^2} \right) \right\}$$

- The second hop: Combine DPC and MAC schemes

- DPC Scheme*

$$R_{p,DPC}^{(2)} = \gamma \left(\frac{\beta P_2}{1 + b^2 \beta P_2} \right)$$

$$R_{c,DPC}^{(2)} = \frac{1}{2} \gamma \left(\frac{(1 - b^2)^2 \bar{\beta}^2 P_2^2}{\sigma_2^4} + \frac{2(1 + b^2) \bar{\beta} P_2}{\sigma_2^2} \right)$$

- MAC Scheme*

$$R_{p,MAC}^{(2)} = \max_{\alpha} \left\{ \min \left[\gamma \left(\frac{b^2 \bar{\alpha} \beta P_2}{\sigma_3^2} \right), \frac{1}{2} \gamma \left(\frac{(1 + b^2) \bar{\alpha} \beta P_2}{\sigma_3^2} \right) \right] + \gamma \left(\frac{\alpha \beta P_2}{1 + b^2 \alpha \beta P_2} \right) \right\}$$

$$R_{c,MAC}^{(2)} = \frac{1}{2} \gamma \left(\frac{(1 + b)^2 \bar{\beta} P_2}{1 + (1 + b^2) \beta P_2} \right)$$

Case II: $a > 1, b > 1$

A natural solution:

- The first hop: the capacity region of the strong interference channel is known as

$$R^{(1)} = \min \left(\gamma(P_1), \frac{1}{2} \gamma((1 + a^2)P_1) \right)$$

- The second hop: the capacity of the Gaussian vector broadcast channel is known as

$$R^{(2)} = \frac{1}{2} \gamma((b^2 - 1)^2 P_2^2 + 2P_2(1 + b^2)).$$

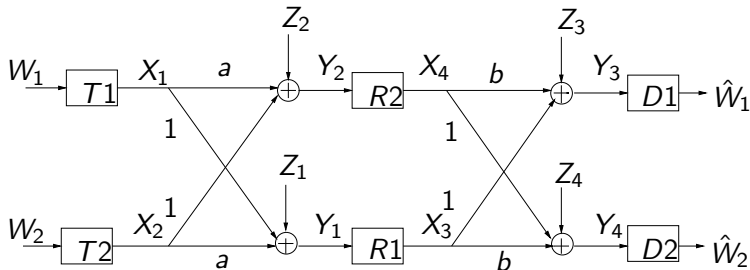
- The overall transmission rate is

$$R = \min\{R^{(1)}, R^{(2)}\}.$$

Optimality in each hop does not guarantee optimality of the entire network!

Case II: $a > 1, b > 1$

After switching the roles of the two relays, the model becomes



- Both hops are converted to weak interference
- We can apply the transmission schemes proposed in Case I
- The resulting rates always include the “natural solution” as a subset

Amplify and Forward

For in-phase relaying, set

$$X_3 = \sqrt{\frac{P_2}{(1+a^2)P_1+1}} Y_1, X_4 = \sqrt{\frac{P_2}{(1+a^2)P_1+1}} Y_2$$

The overall channel becomes

$$\begin{aligned} Y_3 &= (1+ab)X_1 + (a+b)X_2 + Z'_3 \\ Y_4 &= (a+b)X_1 + (1+ab)X_2 + Z'_4 \end{aligned}$$

For out-of-phase relaying, set

$$X_3 = -\sqrt{\frac{P_2}{(1+a^2)P_1+1}} Y_1, X_4 = \sqrt{\frac{P_2}{(1+a^2)P_1+1}} Y_2$$

The overall channel becomes

$$\begin{aligned} Y_3 &= (ab-1)X_1 + (b-a)X_2 + Z'_3 \\ Y_4 &= (a-b)X_1 + (1-ab)X_2 + Z'_4 \end{aligned}$$

Out-of-phase Relaying When $a = b$, $P_1 = P_2$

- When $a = b$, the overall channel becomes two parallel AWGN channel

$$R = \gamma \left(\frac{(1 - a^2)^2 P_1}{1 + a^2 + 1/c^2} \right) = \gamma \left(\frac{(1 - a^2)^2 P_1 P_2}{(1 + a^2)(P_1 + P_2) + 1} \right)$$

- If $a = b < 1$, the sum rate capacity of the interference channel is known at the condition of 'noisy interference'

$$a(a^2 P_1 + 1) \leq \frac{1}{2}, \text{ i.e., } P_1 \leq \frac{1}{a^2} \left(\frac{1}{2a} - 1 \right)$$

and the corresponding symmetric capacity is

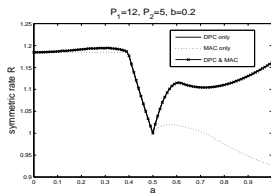
$$C_1 = \gamma \left(\frac{P_1}{1 + a^2 P_1} \right)$$

- The overall rate exceeds the capacity of both hops when

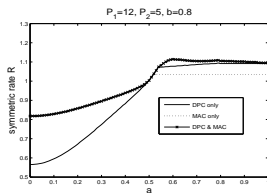
$$\frac{1 + 4a^2 - a^4 + \sqrt{(1 + 4a^2 - a^4)^2 + 4a^2(1 - a^2)^2}}{2a^2(1 - a^2)^2} < P_1 < \frac{1}{a^2} \left(\frac{1}{2a} - 1 \right)$$

For example, when $a = 0.15$, $51.6 < P_1 < 103.7$.

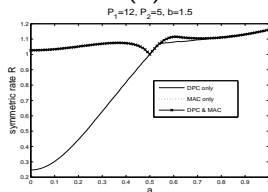
Numerical Examples



(a)



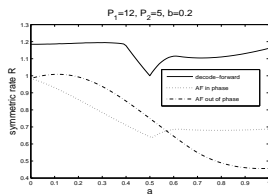
(b)



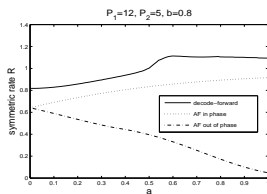
(c)

Figure: Comparison of DPC scheme and MAC scheme in the second hop for the decode-and-forward relaying.

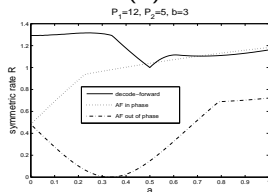
Numerical Examples



(a)



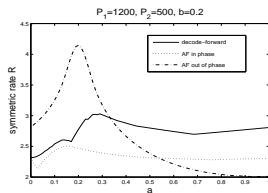
(b)



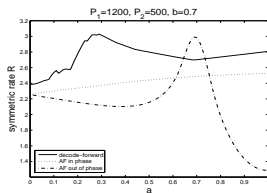
(c)

Figure: Comparison of decode-and-forward relaying and amplify-and-forward relaying in low SNR regime

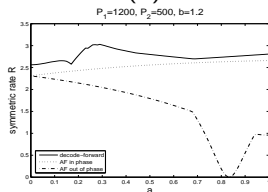
Numerical Examples



(a)



(b)



(c)

Figure: Comparison of decode-and-forward relaying and amplify-and-forward relaying in high SNR regime

Conclusion

- We investigated and compared coding schemes for the two hop interference network under various channel parameters regimes.
- If the first hop has strong interference, i.e., $a > 1$, it is always beneficial to switch the roles of the two relays so that the channel is converted to a weak interference channel with interference gain of $1/a$, and the strength of the second hop is also changed accordingly.
- For the decode-and-forward relaying, the DPC scheme and MAC scheme are both needed for the second hop. Generally however, DPC scheme dominates when b is small and MAC scheme dominates when b is large.
- DF always has better performance than AF except when a is close to b in the high SNR regime.
- When $a = b$, the overall transmission rate can exceed the capacity of both hops.